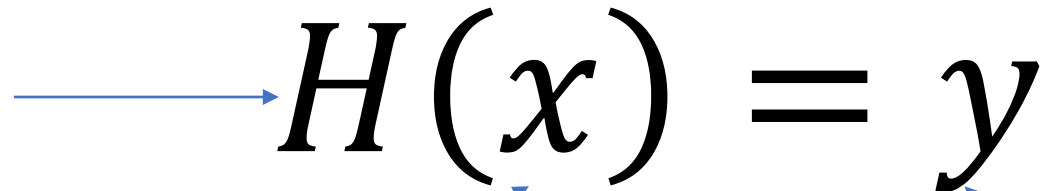


Tutorial -6

SOEN-321

Hash Function

One-way function
There is no H^{-1}



The diagram shows the equation $H(x) = y$. A blue arrow points from the text 'One-way function' to the function symbol H . Another blue arrow points from the text 'Arbitrary length input' to the input variable x . A third blue arrow points from the text 'Short fixed length output' to the output variable y .

$$H(x) = y$$

Arbitrary length input

- Message
- Pre-image

Short fixed length output

- Hash code
- Hash value
- Message digest

Application: Hash Tables with $O(1)$ lookup

Pigeonhole principle

- If n items are put in m containers where $n > m$, then at least one of the containers has more than one item
- It means, there must exist “collision”
 $H(x_1) = H(x_2)$, where $x_1 \neq x_2$



Cryptographic Hash Function

1. Pre-image resistance

- *Given y , it is hard to find x such that $H(x) = y$*

2. Weak collision resistance

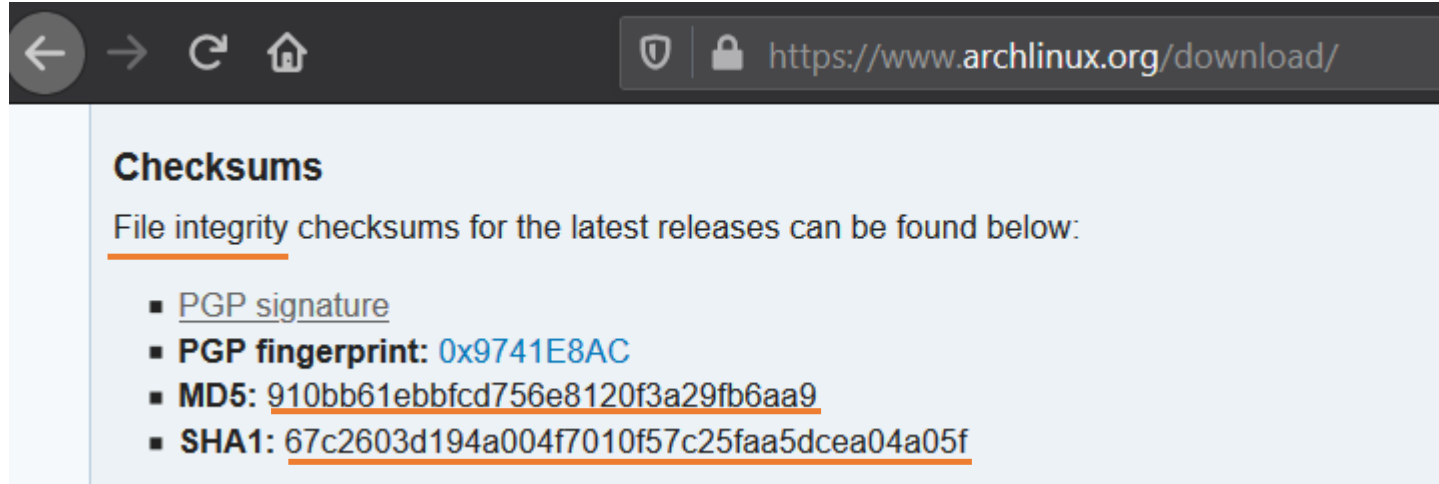
- *Given x_1 , it is hard to find x_2 such that $H(x_1) = H(x_2)$ and $x_2 \neq x_1$*

3. Strong collision resistance

- *it is hard to find any (x_1, x_2) such that $H(x_1) = H(x_2)$*

Examples: MD5, SHA1, SHA2, SHA3

Cryptographic Hash Function



Exercise 5-1

Bob is a paranoid cryptographer who does not trust dedicated hash functions such as SHA1 and SHA-2. Bob decided to build his own hash function based on some ideas from number theory. More precisely, Bob decided to use the following hash function:

$H(m) = m^2 \bmod n$, $n = p \times q$, where p and q are two large distinct primes.

Does this hash function satisfy the one-wayness property? What about collision resistance? Explain.

1- Pre-image resistant:

Yes, since p and q are secret, then finding the square root *mod* n is a hard problem

2-Weak collision resistant

No, since for any given input m , the attacker can get the same hash value using input $-m$

3-Strong collision resistant

No, it is easy to choose any pair $(m, -m)$ which yields the same hash

Exercise 6-1

Let $x=111$ and $y=19301$. Factor $n=21311$ using the fact that $x^2 \equiv y^2 \pmod n$.

$$x^2 - y^2 \equiv 0 \pmod n$$

$$(x + y)(x - y) \equiv 0 \pmod n$$

If n divides $(x + y)(x - y)$, then they share common factors

$$\gcd(x \pm y, n) = p \text{ or } q$$

$$\gcd(111 + 19301, 21311)$$

$$\gcd(19412, 21311)$$

$$21311 = 19412 + 1899$$

$$19412 = 10 \times 1899 + 422$$

$$1899 = 4 \times 422 + 211$$

$$422 = 2 \times 211 + 0$$

$$\gcd(19412, 21311) = p = 211$$

$$q = \frac{n}{p} = \frac{21311}{211} = 101$$